

HW#2, Assigned Sept 2, Due Sept 15.

1) *Network Information Theory*. Suppose a Markov chain starts in one of n values, i.e. "states", necks down to one of $k < n$ values and then fans back to $m > k$ states. That is

$X \rightarrow Y \rightarrow Z$ is a Markov Chain in that order and

$$X \in \{1, 2, \dots, n\}, Y \in \{1, 2, \dots, k\}, Z \in \{1, 2, \dots, m\}$$

(a) Show that the dependence of X and Z is limited by the bottleneck by proving that $I(X; Z) \leq \log k$:

(b) Evaluate $I(X; Z)$ for $k = 1$, and conclude that no dependence can survive such a bottleneck.

2) *Data Storage Information Theory*. The following sequences of bits are very common in data storage devices.

Let $X = (X_1, X_2, \dots, X_n)$ be (possibly dependent) binary random variables. Suppose one calculates the lengths of runs of zeros or ones and let $R = (R_1, R_2, \dots, R_m)$ be a sequence of runs. For example, the sequence $X = (0001100100)$ yields run lengths $R = (3; 2; 2; 1; 2)$. As a note, in CD and DVD any run must satisfy $3 \leq R_i \leq 11$ for all i .

Compare $H(X_1, X_2, \dots, X_n), H(R), H(X_n, R)$, namely find as many inequality relations or expressions relating each.

3. Consider a sequence of n binary random variables (X_1, X_2, \dots, X_n) . Each sequence with an even number of 1's has probability $2^{-(n+1)}$ and each sequence with an odd number of 1's has probability 0. Find the mutual informations,

$$I(X_1; X_2), I(X_2; X_3 | X_1), \dots, I(X_{n-1}; X_n | X_1, \dots, X_{n-2}),$$

4. Let each term in (X_1, X_2, \dots, X_n) be i.i.d. drawn according to probability mass function $p(x)$. Find

$$\lim_{n \rightarrow \infty} [p(X_1, X_2, \dots, X_n)]^{1/n}$$

5. The problem relates both to the AEP and decision theory where given a set of data we are trying to find which distribution it comes from.

Let X_1, X_2, X_3, \dots be independent identically distributed random variables drawn according to the probability mass function $p(x), x \in \{1, 2, \dots, m\}$ $p(x)$.

Thus $p(x_1, \dots, x_n) = \prod_{i=1}^n p(x_i)$. We know that $-\frac{1}{n} \log p(x_1, \dots, x_n) \rightarrow H(X)$. Let

$q(x_1, \dots, x_n) = \prod_{i=1}^n q(x_i)$ where $q(x)$ is another probability mass function on $\{1, 2, \dots, m\}$

(a) Evaluate $\lim_{n \rightarrow \infty} -\frac{1}{n} \log q(x_1, \dots, x_n)$ where X_1, X_2, X_3, \dots are i.i.d. each X_1, X_2, X_3, \dots , with density $p(x)$

(b) Now evaluate the limit of the log likelihood ratio

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log \frac{q(x_1, \dots, x_n)}{p(x_1, \dots, x_n)}$$

when X_1, X_2, X_3, \dots are iid with probability mass function $p(x)$. This result shows that the odds favoring $q(x)$ are exponentially small when $p(x)$ is true.