

Breaking Smart Card Based ElGamal Signature and Its Variants

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Yuliang Zheng
Monash University

Tsutomu Matsumoto
Yokohama Nat. University

Exploiting hardware faults



- **Secret key algorithms**

- DES, IDEA, FEAL, etc --- broken by Biham & Shamir (18 Oct. & 30 Oct.)
- (see also J-J Quisquater, 23 Oct.)

- **Public key signature**

- RSA --- broken by
 - ☉ Bellcore (25 Sept., no details were published)
 - ☉ Nat. Uni. of Singapore (23 Oct.)
- ***ElGamal family -- ???***

Completing the big picture



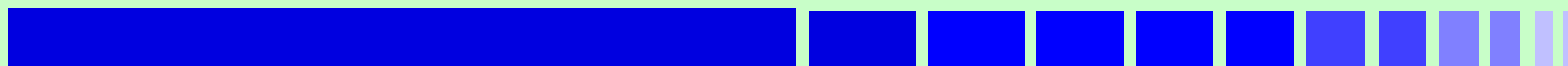
- **We show how to break smart card based**
 - **ElGamal signature and**
 - **all its variants,**
by (temporally) falsifying part of the
circuitry in the smart card
- **Time: 31 Oct. 1996**

The Attack Model



- **An attacker, who is in possession of a target smart card, may introduce faults into tamper-proof hardware in the smart card, say by exposing it to certain physical effects (heat, laser, pressure, radiation, etc)**

The Attack Model (cnt'd)



- **The attack may then compromise a secret in the smart card.**
 - **a DES encryption / decryption key**
 - **a RSA signature-generation / decryption key**
 - **certain authentication / identification cards**
 - **.....**

ElGamal Signature

- **Alice's keys**

- x_a --- secret key for signature generation

- $y_a = g^{x_a} \bmod p$ --- public key

- **Alice's signature on a message m is a pair of number (r,s) :**

- $x \in_R [1, \dots, p-1]$
 $r = g^x \bmod p$

- $s = \frac{\text{hash}(m) - x_a \cdot r}{x} \bmod (p-1)$

Idea for finding x_a



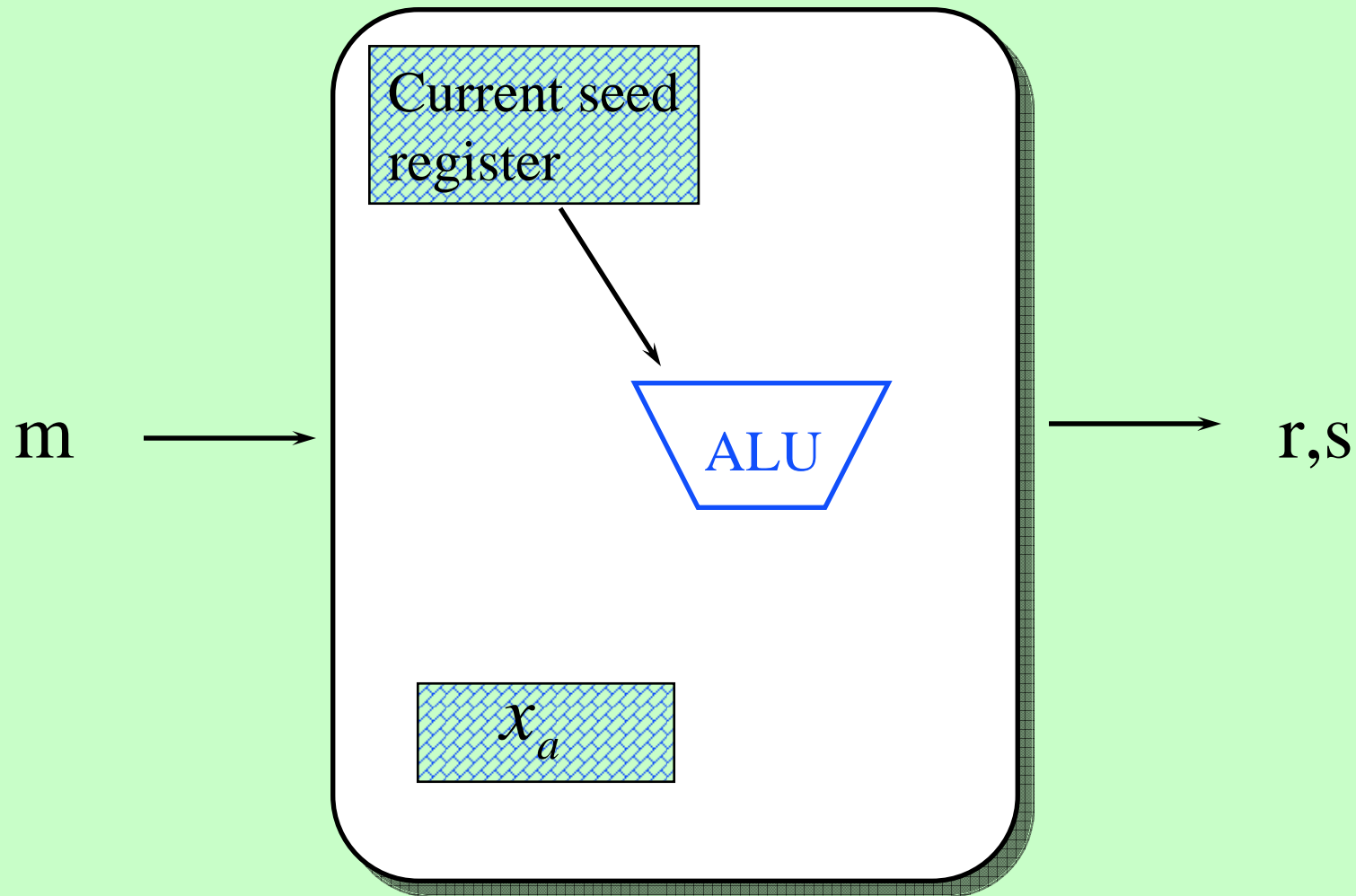
- An attacker may figure out x_a by making the “random” number x predictable, through the falsification of part of the hardware
 - software pseudo-random number generator --- suppressing the contents of the status register that stores the current “seed” number
 - hard pseudo-random number generator --- suppressing its output

Two possibilities



- **software pseudo-random number generator**
 - suppressing the contents of the status register that stores the current “seed” number
- **hard pseudo-random number generator**
 - suppressing its output

When software PRG is used

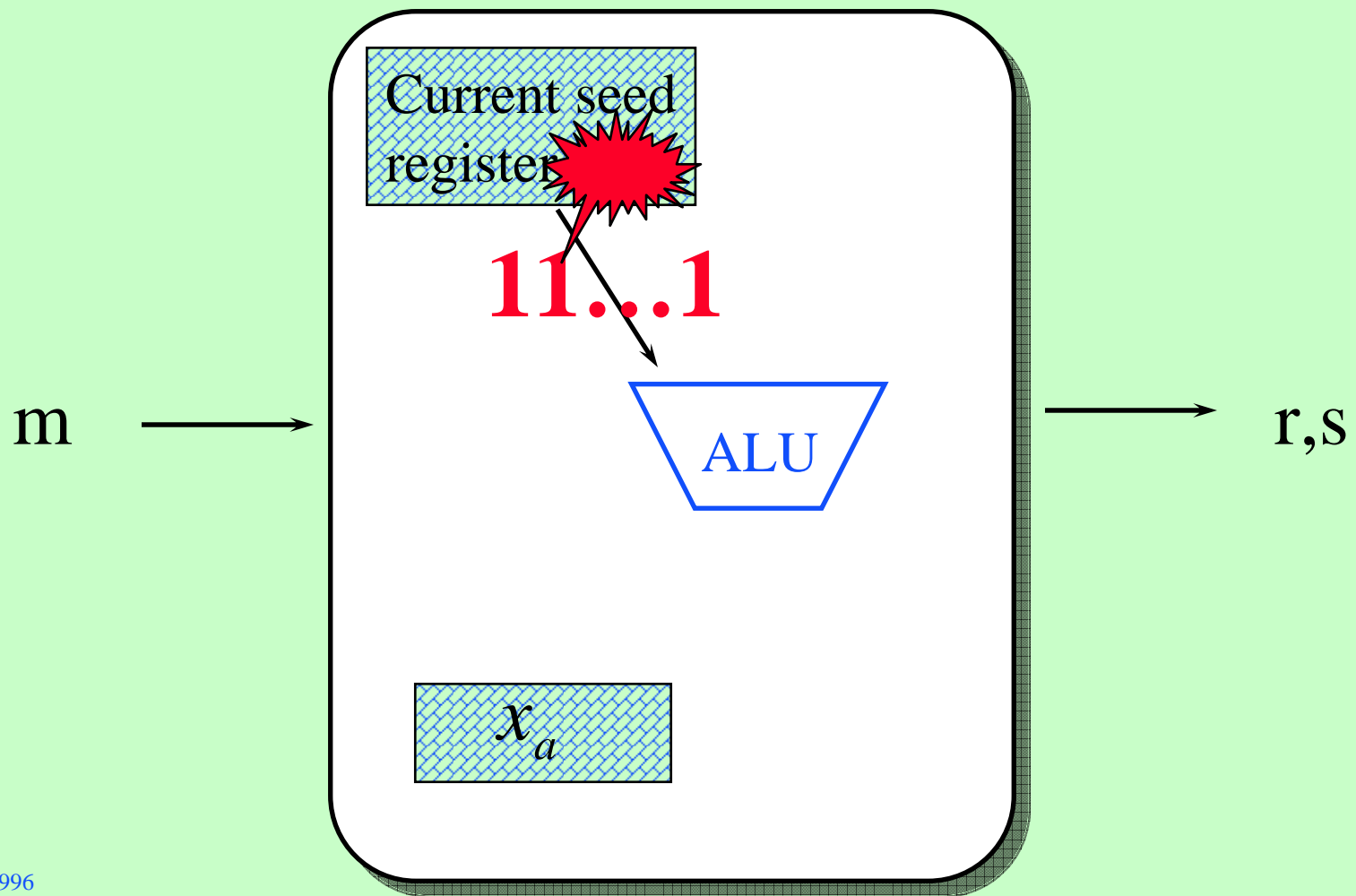


Main idea

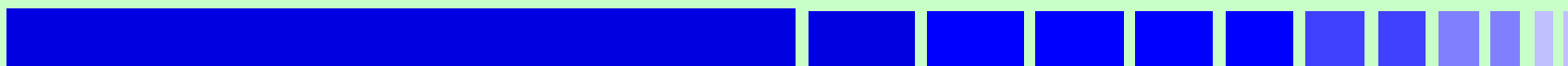


- **suppress the contents of the status register that stores the current seed for the PRG**
- **so that only a fixed number, such as the all-1 value, is (temporarily) available to the CPU**
- **collect the output (r,s) of the smart card**

When software PRG is used



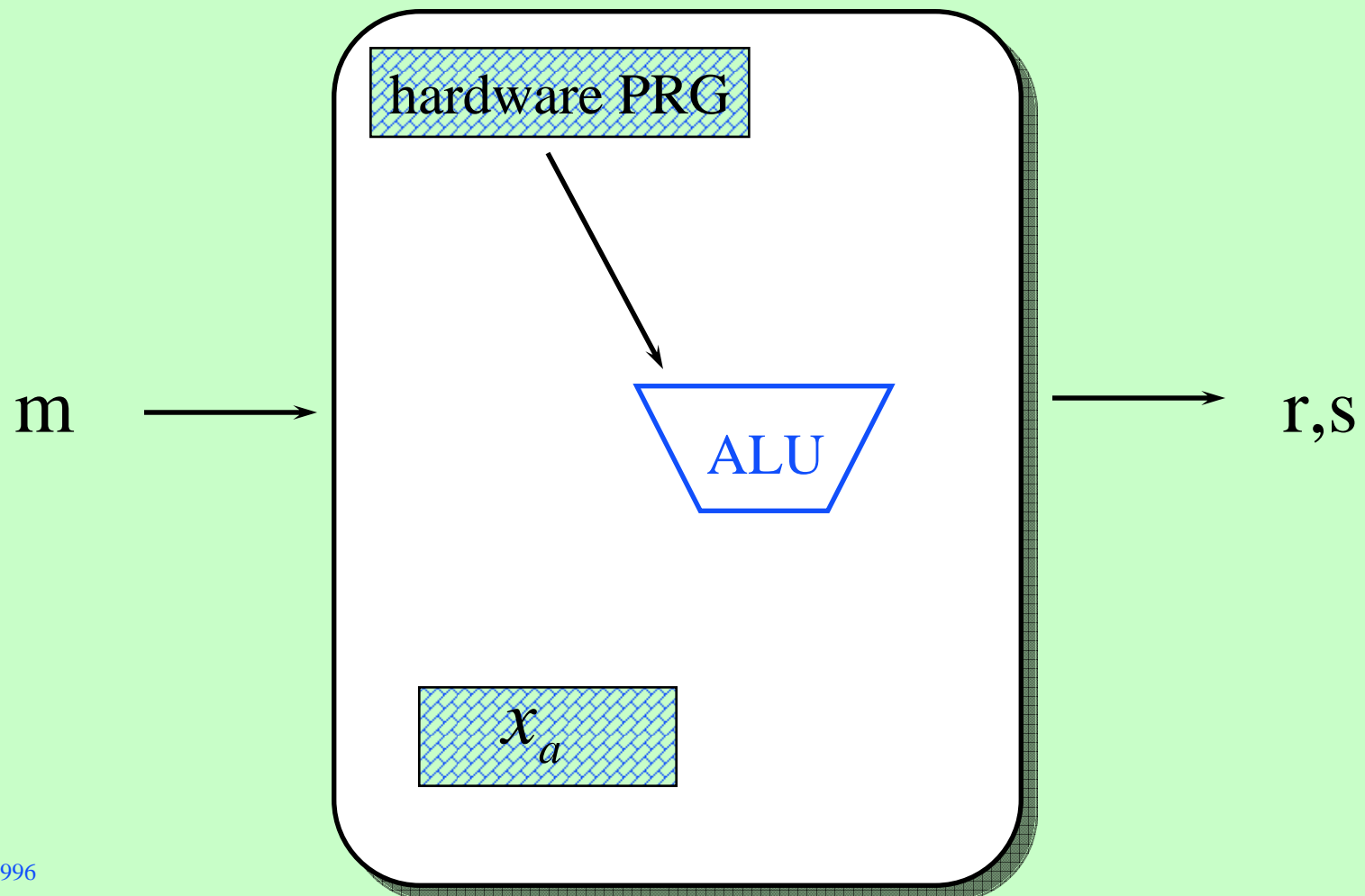
Obtain the secret key x_a



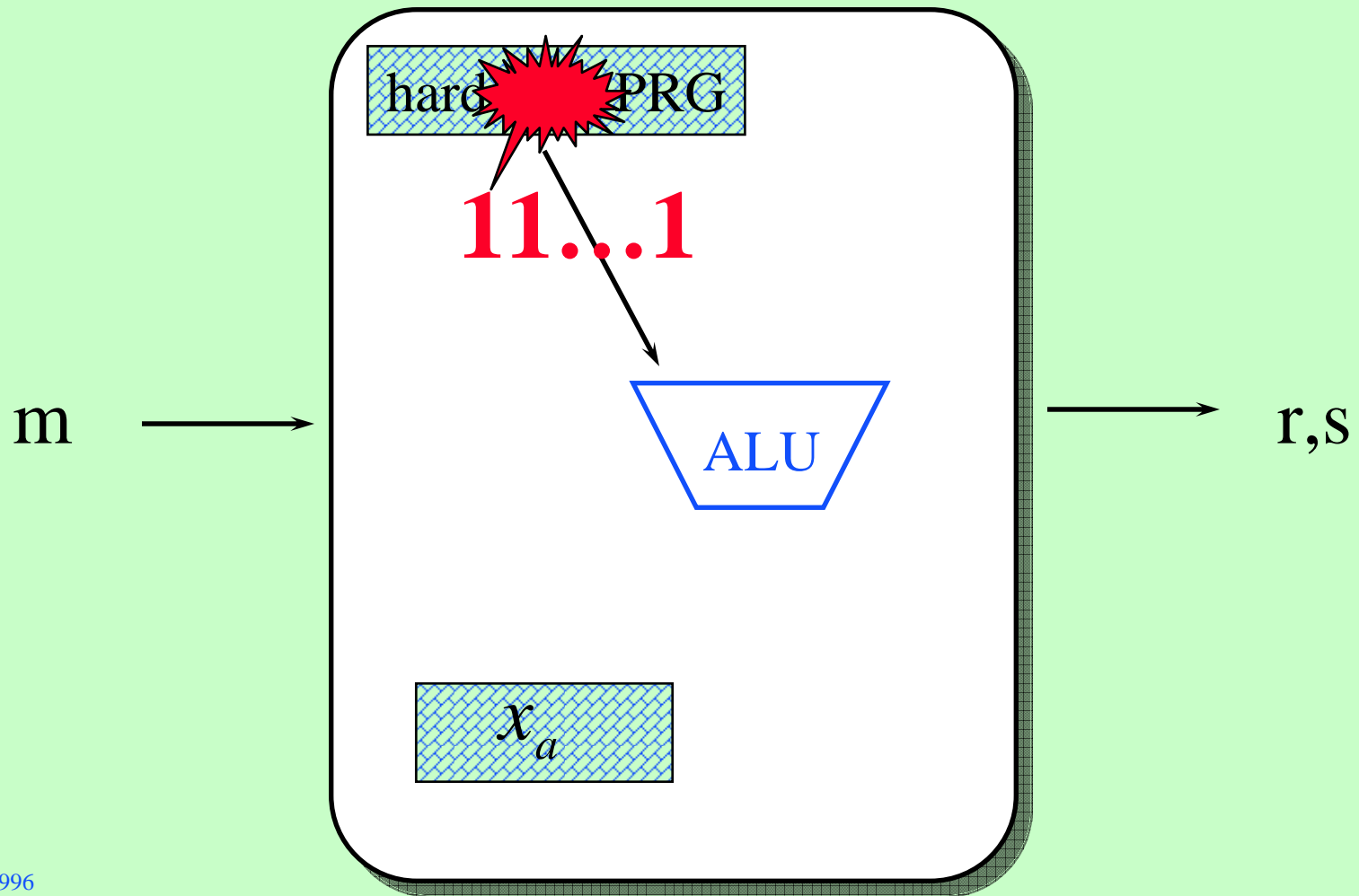
- The attacker can now calculate x from
 - the fixed seed, say the all-1 value
 - the (public) PRG algorithm
- He can then find out x_a

$$x_a = \frac{\mathit{hash}(m) - s \cdot x}{r} \bmod (p - 1)$$

When hardware PRG is used



Suppressing output of hardware PRG



Obtain the secret key x_a



- the attacker computes

$$x_a = \frac{\text{hash}(m) - s \cdot x}{r} \bmod (p - 1)$$

Another attack



- Some smart cards have a built-in PRG that would produce a constant output when a lower than normal voltage is supplied.
- For such a smart card, an attacker may obtain Alice's secret key x_a by asking the card to sign 2 different messages while keeping the voltage low.

Another attack (cnt'd)

- We will have

$$s_1 = \frac{\text{hash}(m_1) - x_a \cdot r}{x} \bmod (p-1)$$

$$s_2 = \frac{\text{hash}(m_2) - x_a \cdot r}{x} \bmod (p-1)$$

- Hence

$$\square \quad x = \frac{\text{hash}(m_1) - \text{hash}(m_2)}{s_1 - s_2} \bmod p-1$$

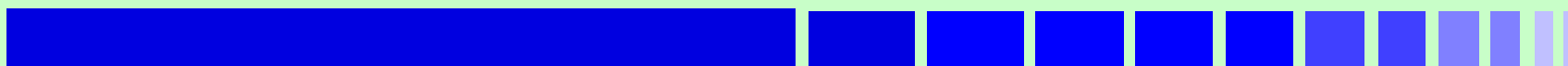
$$x_a = \frac{\text{hash}(m_1) - x \cdot s_1}{r} \bmod p-1$$

Breaking other ElGamal variants



- **All other ElGamal variants can be broken in a similar way**
 - **DSS**
 - **Schnorr**
 - **.....**
 - **(Elliptic curve based schemes)**

Attacking RSA by NUS



- Nat Uni Singapore method
 - Finding an RSA decryption / signature generation key by

complementing *one or a few* bits at random positions in the unknown decryption/signature-generation key,

which is harder than suppressing the entire contents of a register !

A comparison with RSA (cnt'd)



- Therefore, compared to the attacker on RSA signature / decryption card, our attack seems
 - simpler
 - more feasible to mount
 - ☺ say, on an entire data-bus
 - but, harder to prevent
 - ☺ reason --- (r,s) is a perfectly valid signature !

Smart cards under threat

--- the big picture ---

