

Texas Instruments MSP430FR6989 Launchpad

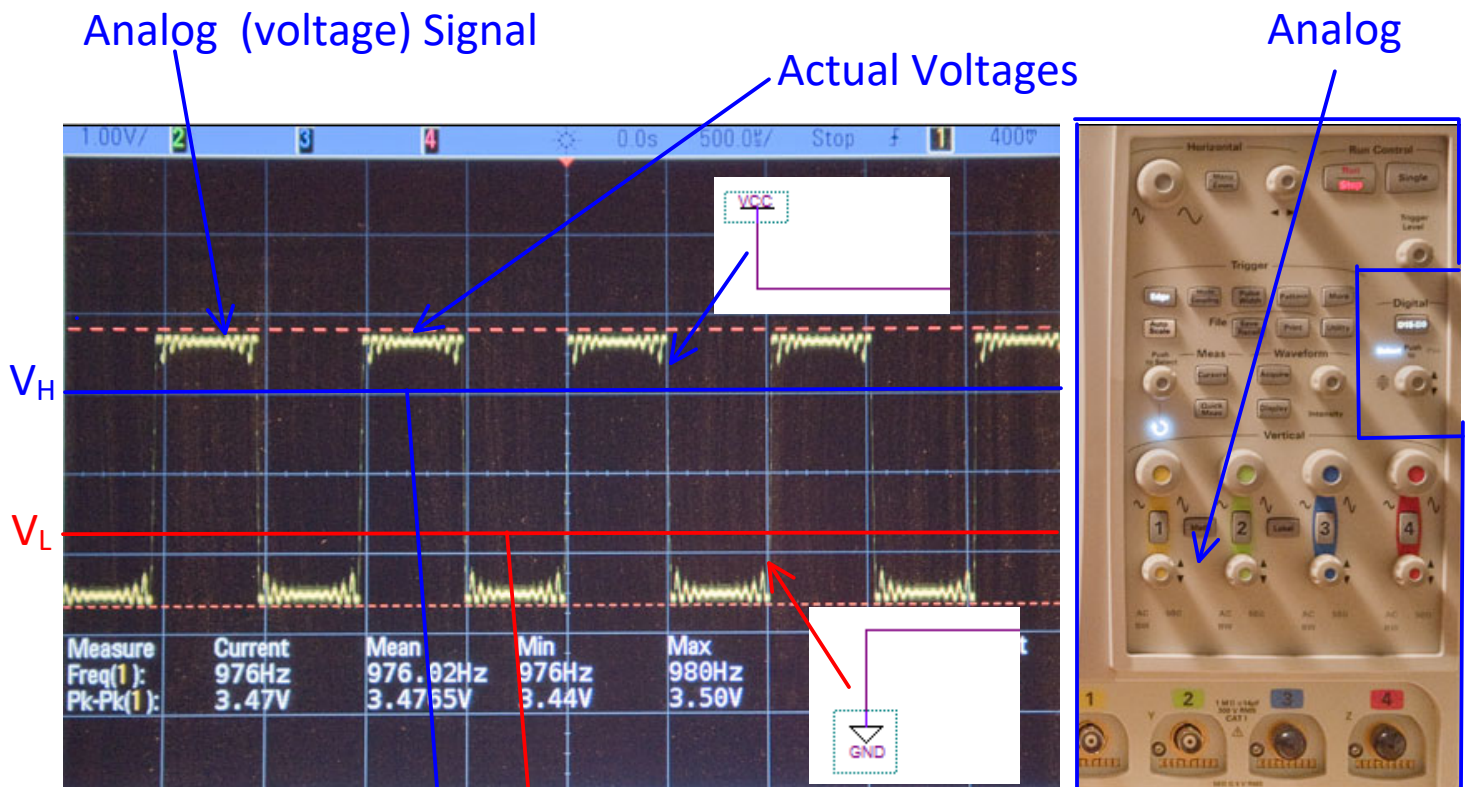


MCU:
MSP430FR6989IPZ

MSP-EXP430FR6989 LaunchPad Development Kit

MSP430FR6989 Microcontroller

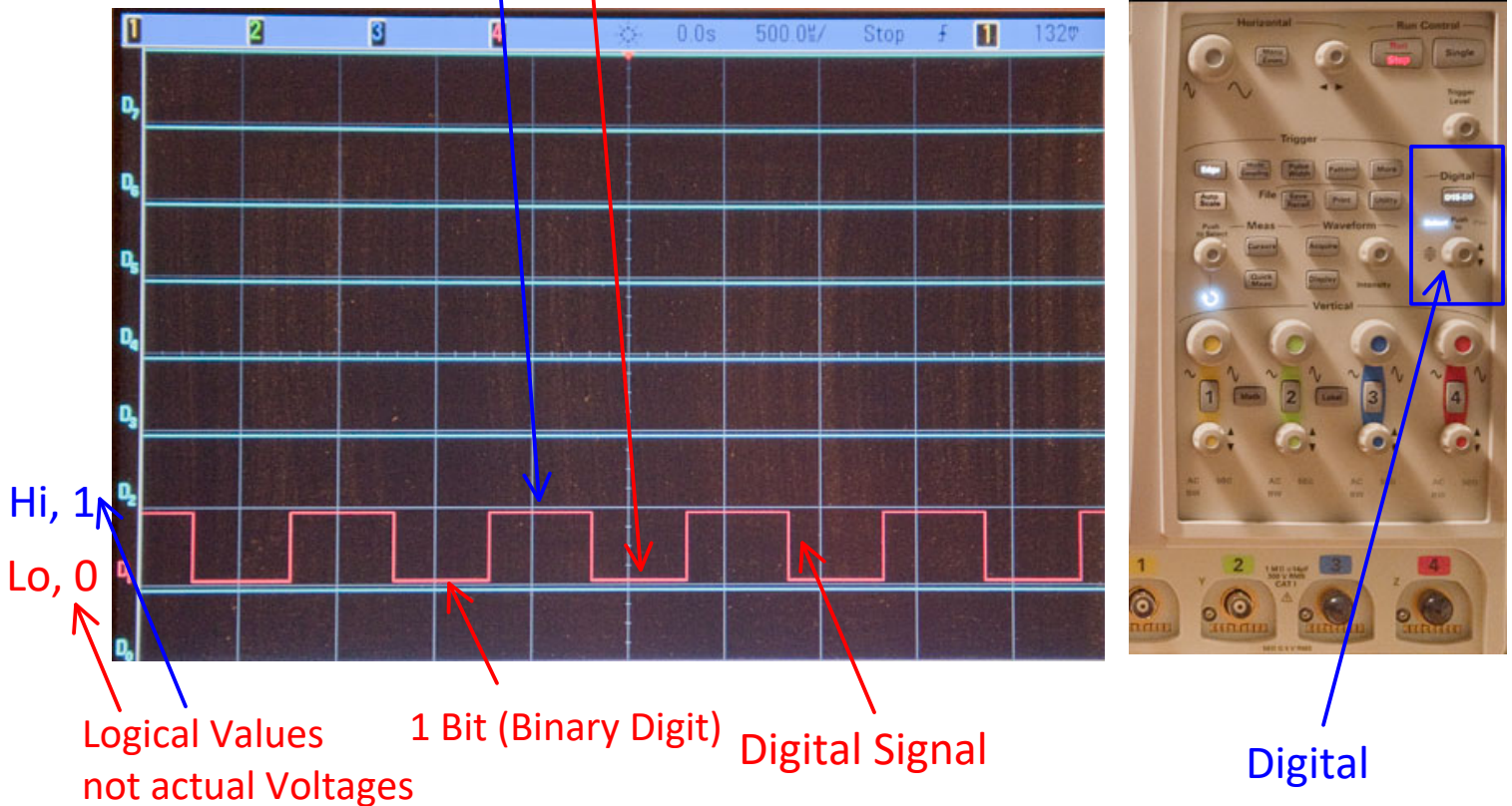
- 1.8-V to 3.6-V operation
- 16-bit RISC architecture up to 16-MHz system clock and 8-MHz FRAM access
- 128KB of nonvolatile FRAM
- 83 GPIOs



Interpretation

Digital signal is Hi or 1 if analog signal voltage > blue voltage

Digital signal is Lo or 0 if analog signal voltage < red voltage



Number Systems

Binary numbers
base, radix = 2

$A_{n-1} \quad A_1 A_0 \quad A_{-1} A_{-2} \quad A_{-m}$

$A_i \in \{0, 1\}$ $A_i \rightarrow$ binary digit (bit)

Binary number

$\rightarrow 1011011101 = (1011011101)_2$
 bit bit binary point

Converting from binary to decimal

$$(11010111)_2 = 1 \times 2^5 + 1 \times 2^4 + 0 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$$

$$\begin{array}{cccccccc} 5 & 4 & 3 & 2 & 1 & 0 & -1 & -2 \\ 1 & 1 & 0 & 1 & 0 & 1 & 1 & 1 \end{array}$$

$$+ 1 \times 2^{-1} + 1 \times 2^{-2} = (53.75)_{10}$$

Notation: Code Composer Studio

$(110101)_2 \rightarrow 110101b$ binary

msd (most significant digit) $\rightarrow$ 1 1 0 1 0 1 $\leftarrow$ lsd (least significant digit)

decimal integer to binary

$$(41)_{10} = (?)_2$$

$$\begin{array}{rcl}
 \frac{41}{2} & = & 20 + \frac{1}{2} \\
 \frac{20}{2} & = & 10 + \frac{0}{2} \\
 \frac{10}{2} & = & 5 + \frac{0}{2} \\
 \frac{5}{2} & = & 2 + \frac{1}{2} \\
 \frac{2}{2} & = & 1 + \frac{0}{2} \\
 \frac{1}{2} & = & 0 + \frac{1}{2}
 \end{array}$$

$(101001)_2$

Hexadecimal numbers

Hexadecimal radix = 16 $16 = 2^4$ radix is a power of 2

useful to represent binary numbers easily in compact form (easier on human beings than binary representation)

0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F

¹⁰ ¹¹ ¹² ¹³ ¹⁴ ¹⁵
 ↗ ↗ ↗ ↗ ↗ ↗

16 hexadecimal digits

$$(B65F)_{16} = 11 \times 16^3 + 6 \times 16^2 + 5 \times 16^1 + 15 \times 16^0$$

$$= (46687)_{10}$$

In Code Composer Studio :

$(B65F)_{16} \rightarrow \text{0XB65F}$ (C style)

$\rightarrow \text{B65FH}$ (TI style)

Important table for converting bases among binary, octal and hexadecimal number representation

Decimal	Binary	Octal	Hex
00	0000	00	0
01	0001	01	1
02	0010	02	2
03	0011	03	3
04	0100	04	4
05	0101	05	5
06	0110	06	6
07	0111	07	7
08	1000	10	8
09	1001	11	9
10	1010	12	A
11	1011	13	B
12	1100	14	C
13	1101	15	D
14	1110	16	E
15	1111	17	F

Converting from binary to hexadecimal

$$(10110001101011)_2 = (?)_{16}$$

(use Table)

$$\begin{array}{cccc} \overbrace{0010}^{\text{2}} & \overbrace{1100}^{\text{C}} & \overbrace{0110}^{\text{6}} & \overbrace{1011}^{\text{B}} \end{array}$$

pad zeros if required start from here

$$= (2C6B)_{16}$$

Converting from hexadecimal to binary

$$(3A6)_{16} = (?)_2$$

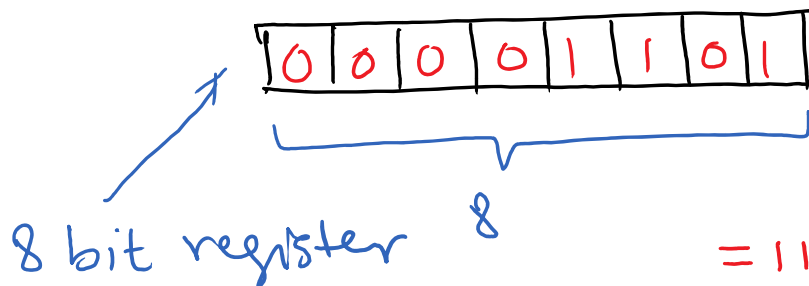
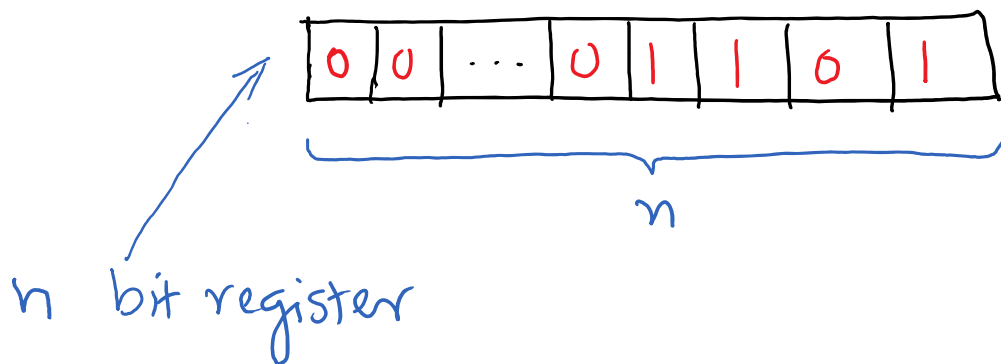
(use Table)

$$\begin{array}{cccc} (3A6)_{16} = & \overbrace{0011}^{\text{discard}} & \overbrace{1010} & \overbrace{0110} \end{array}$$

$$= (1110100110)_2$$

Complements of Binary numbers.

Binary Number $N = 1101$ Stored in an n bit register.



$= 1101$ stored in an 8 bit register

$$N = 1101$$

$$n = 8 \text{ (word length)}$$

1's Complement of Binary Numbers

$$1's \text{ complement of } N \rightarrow (2^n - 1) - N$$

register size
number

Change all the 1's to 0, and all the 0's to 1 in N to get 1's complement of N

$$\begin{array}{ccc}
 10110010 & \xrightarrow{1's \text{ comp.}} & 01001101 \\
 111 & \xrightarrow{1's \text{ comp.}} & 11111000
 \end{array}$$

$$2's \text{ complement of } N \rightarrow \begin{cases} 2^n - N & \text{for } N \neq 0 \\ 0 & \text{for } N = 0 \end{cases}$$

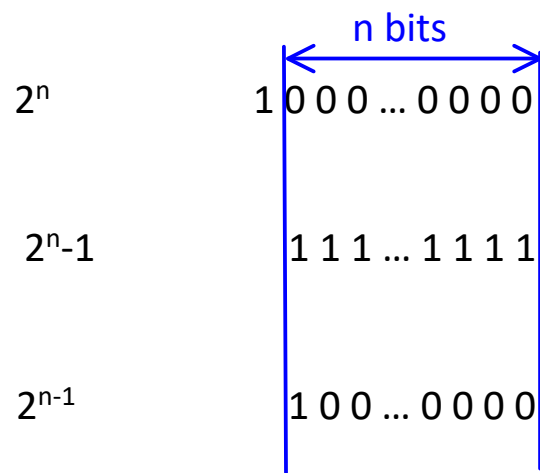
$$2's \text{ comp. of } N = 1's \text{ comp. of } N + 1$$

$$\begin{array}{rcl}
 101100 & \xrightarrow{2's \text{ comp}} & 010011 \leftarrow 1's \text{ comp} \\
 & & + 1 \\
 \hline
 & & 010100 \leftarrow 2's \text{ comp.}
 \end{array}$$

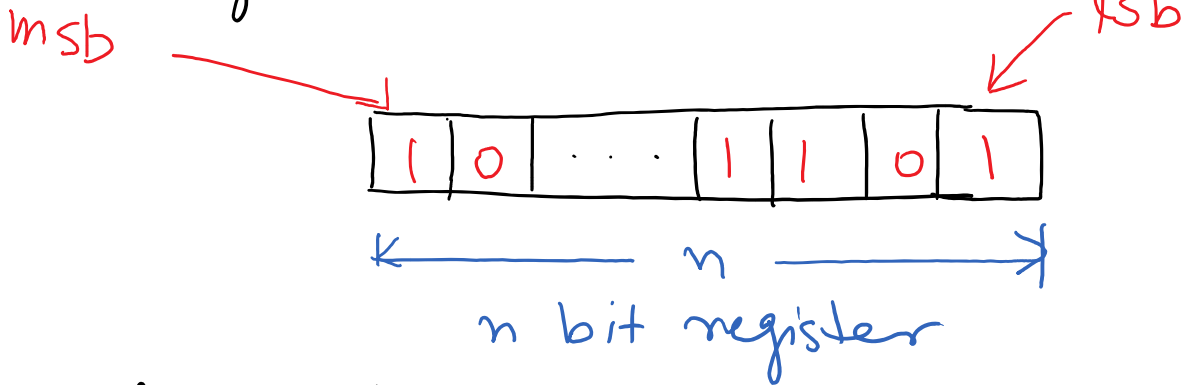
"Trick" to find 2's complement of N:

Leave all the least significant zeros and the first 1 unchanged - replace all the remaining 0's by 1's, and 1's by 0's.

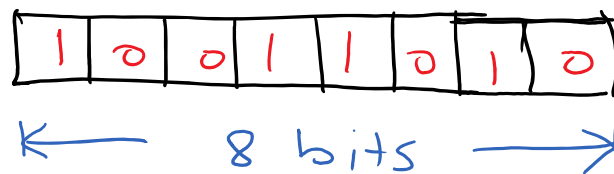
$$2's \text{ comp. of } \underbrace{101100}_{\substack{\text{invert} \\ \text{these}}} \longrightarrow \underbrace{010100}_{\substack{\text{leave these} \\ \text{unchanged}}}$$



Unsigned Numbers

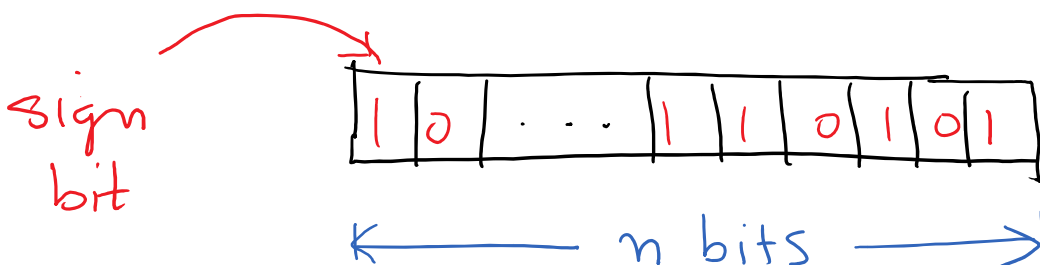


All the bits in the register represent the magnitude of the number.



$$\text{Number } N = +(10011010)_2$$

Signed Numbers (2's complement representation)



example 8 bit registers

$$N = (9)_{10} \longrightarrow \begin{array}{ccccccc} 0 & 0 & 0 & 0 & | & 0 & 0 & 1 \end{array}$$

Sign bit
0 = +
magnitude
= 0001001 = (9)₁₀

$$N = (-9)_{10} \longrightarrow \text{2's comp of } 00001001$$

$$= 11110111$$

Sign bit
1 = -

Example

$$\begin{array}{ccccccc} 1 & 1 & 1 & 0 & 1 & 0 & 1 & 0 \end{array}$$

signed bit is 1 $\therefore$ number is -ve

2's comp. of 11101010 is 00010110 = (22)₁₀

$\therefore$ 11101010 is (-22)₁₀

$$\begin{array}{ccccccc} 0 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \end{array}$$

signed bit 0 $\therefore$ number is +ve

00011011 is (27)₁₀

Signed Binary Addition and subtraction

Addition

Add the two numbers (including the signed bit)
discard the carry out

$$\begin{array}{r}
 X = +6 \quad 00000110 \\
 Y = +13 \quad 00001101 \\
 \hline
 Z = X + Y = +19 \quad 00010011 \\
 \quad \quad \quad \swarrow \text{signbit}
 \end{array}$$

$$\begin{array}{r}
 X = -6 \quad 11111010 \\
 Y = +13 \quad 00001101 \\
 \hline
 Z = X + Y = +7 \quad 10000111 \\
 \quad \quad \quad \text{discard } \textcircled{1} \quad \swarrow \text{signbit}
 \end{array}$$

$$\begin{array}{r}
 X = +6 \quad 00000110 \\
 Y = -13 \quad 11110011 \\
 \hline
 Z = X + Y = -7 \quad 11110001 \\
 \quad \quad \quad \swarrow \text{signbit}
 \end{array}$$

2's comp of 11111001 is 00000111 = +7
therefore 11111001 is -7



$$\begin{array}{r}
 X = -6 \quad 11111010 \\
 Y = -13 \quad 11110011 \\
 \hline
 Z = X + Y = -19 \quad \text{①} 11101101 \\
 \text{discard} \quad \text{sign bit}
 \end{array}$$

2's comp of 11101101 is $00010011 = +19$
 therefore 11101101 is -19

Subtraction (2's comp. representation)

$$Z = X - Y$$

$$Z = \left\{ \begin{array}{l} X + (\text{2's comp of } Y) \\ \text{discard the carry out} \end{array} \right.$$

$$\begin{array}{r}
 X = -6 \quad 11111010 \\
 Y = -13 \xrightarrow{\text{2's comp of } -13} 00001101 \quad \left. \vphantom{\begin{array}{l} X \\ Y \end{array}} \right\} \text{add} \\
 \hline
 Z = X - Y = +7 \quad \text{①} 00000111 \\
 -6 - (-13) \quad \text{discard} \quad \text{sign bit}
 \end{array}$$

$$\begin{array}{r}
 X = +6 \quad 00000110 \\
 Y = -13 \xrightarrow{\text{2's comp of } -13} 00001101 \quad \left. \vphantom{\begin{array}{l} X \\ Y \end{array}} \right\} \text{add} \\
 \hline
 Z = X - Y = +19 \quad 00010011 \\
 +6 - (-13) \quad \text{sign bit}
 \end{array}$$



overflow in 2's complement representation

Assume 8 bits

Sign bit
 $0 \underbrace{1111111}_{127} \rightarrow +127 \leftarrow \text{maximum positive number}$

$1111111 \rightarrow -1$

$10000001 \rightarrow -127$

$10000000 \rightarrow -128 \leftarrow \text{maximum negative number}$

complementing the max. negative number $\rightarrow$ overflow
 Note: 2's complement of 10000000 should give $+128$

but $+128$ needs more than 8 bits!

$$\begin{array}{r} \text{9 bits} \\ 010000000 \\ + \end{array} \rightarrow +128$$



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keep this "special case" in mind when converting
the maximum negative number

8 bits can store a number if $-128 \leq \text{number} \leq +127$

otherwise there will be overflow

Range of 2's Comp

$$-2^{n-1} \leq N \leq 2^{n-1} - 1$$



Addition overflow in 2's complement representation

8 bits 2's complement representation

$-128 \leq N \leq 127$

Discard

addition

10000001 = -127

10000001 = -127

100000010 = 2 overflow!!

(2's complement of 10000001 is 01111111 = 127
therefore 10000001 = -127)

11000001 = -63

+ 11000001 = -63

-63 + (-63) = -126

110000010 = -126 no overflow!!

(2's complement of 10000010 is 01111110 = 126
therefore 10000010 = -126)

Carry out of signed bit position **not equal to**

Carry into signed bit position

↓

1

10

10000001 = -127

+ 10000001 = -127

100000010 = 2 overflow!!

Carry out of signed bit position **is equal to**
 Carry into signed bit position

11

11000001 = -63

+ 11000001 = -63

110000010 = -126 no overflow!!

Discard

For addition (in 2's compliment representation),

- Overflow can only happen if both numbers have the same sign
- **There is overflow** if carry out of the signed bit position **is not equal to** carry into the signed bit position