

Math 3700 11/15 and 11/17

Today: other isom thms

HW 9

- ↳ SES
- ↳ gp extensions
- ↳ free gps
- ↳ gp presentations

Isomorphism Thms at Lightning Speed

① Let $\phi: G \rightarrow G'$ be a gp hom. Then

$$G/\ker\phi \cong G'$$

② Let $H \trianglelefteq G$ and $N \trianglelefteq G$. Then

$$\begin{array}{c} HN = \{hn \mid h \in H, n \in N\} \\ \trianglelefteq \\ N \trianglelefteq H \\ \Downarrow \\ N \cap H \trianglelefteq H \end{array}$$

and $HN/N \cong H/N \cap H$.

③ Let $N \trianglelefteq G$. Then

$$+ ④ \left\{ \begin{array}{l} \text{normal} \\ \text{subgps} \\ \text{of } G/N \end{array} \right\} \leftrightarrow \left\{ \begin{array}{l} \text{normal} \\ \text{subgps of } G \\ \text{containing } N \end{array} \right\}$$

$$M/N \quad \text{and} \quad (G/N)/(M/N) \cong G/M.$$

Note: $|G/M| = |G/N| \cdot |M/N|$.

Examples using them

last time

Prove: $\text{lcm}(a,b)/\text{gcd}(a,b) = ab$

$$\begin{aligned} \text{Let } G = \mathbb{Z}, H = a\mathbb{Z}, N = b\mathbb{Z}. \text{ Then} \\ a\mathbb{Z} + b\mathbb{Z} &= \text{gcd}(a,b)\mathbb{Z} & ax + by \in \text{gcd}(a,b)\mathbb{Z} \\ a\mathbb{Z} \cap b\mathbb{Z} &= \text{lcm}(a,b)\mathbb{Z} & \forall x, y \in \mathbb{Z} \\ \text{so } \text{gcd}\mathbb{Z}/b\mathbb{Z} &\cong a\mathbb{Z}/\text{lcm}\mathbb{Z} \Rightarrow \frac{\text{gcd}}{b} = \frac{a}{\text{lcm}}. \end{aligned}$$

Recall $\text{SL}_n(\mathbb{R}) \trianglelefteq \text{GL}_n(\mathbb{R})$ and $\text{SL}_n(\mathbb{R}) = \ker(\det)$
for $\det: \text{GL}_n(\mathbb{R}) \rightarrow \mathbb{R}^\times$. By 1st iso, $\text{GL}_n(\mathbb{R})/\text{SL}_n(\mathbb{R}) \cong \mathbb{R}^\times$
so 3rd iso: $\left\{ \begin{array}{l} \text{subgps} \\ \text{of } \mathbb{R}^\times \end{array} \right\} \leftrightarrow \left\{ \begin{array}{l} \text{subgps of } \text{GL}_n(\mathbb{R}) \\ \text{containing } \text{SL}_n(\mathbb{R}) \end{array} \right\}$

Since all subgps of \mathbb{R}^\times are normal, all subgps of $\text{GL}_n(\mathbb{R})$ containing $\text{SL}_n(\mathbb{R})$ are normal (!)

Short Exact Sequences of Groups

A short exact sequence (SES) of groups is notation

$$e \rightarrow N \xrightarrow{i} G \xrightarrow{p} Q \rightarrow e' \quad \text{often 1 or 0 (if Abelian)}$$

which neatly packages a lot of info:

- N, G, Q are groups and i, p are gp homs
- i is injective $1 \rightarrow N \xrightarrow{i} G$
- p is surjective $G \xrightarrow{p} Q \rightarrow 1$
- $\ker(p) = \text{im}(i)$ $N \xrightarrow{i} G \xrightarrow{p} Q$

What is this telling us?

① We know $\ker(p) \cong G$, and by 1st iso

$$G/\ker(p) \cong \text{im}(p)$$

② But p is surjective, so $G/\ker(p) \cong \text{im}(p) = Q$.

③ By exactness, $\ker(p) = \text{im}(i)$ so $G/\text{im}(i) \cong Q$.

④ But i is injective, so $N \cong \text{im}(i)$, hence

$$"G/N \cong Q"$$

The gp G is called the extension of Q by N . ↪ quotations b/c its really $G/i(N) \cong Q$ but $i(N) \cong N$ sooo ...

Sometimes, $G \cong N \times G/N \cong N \times Q$. ↪ not always!

Examples (from HW9 Exc1)

1. $0 \rightarrow 2\mathbb{Z} \xrightarrow{i} \mathbb{Z} \xrightarrow{q} \mathbb{Z}/2 \rightarrow 0$
- | | |
|-----------------------|---|
| $a \mapsto a$ | (i) everything involved is a gp + gp hom |
| $b \mapsto b \bmod 2$ | (ii) i is injective since if $a, b \in 2\mathbb{Z}$ are equal in \mathbb{Z} , then equal in $2\mathbb{Z}$. |
| | (iii) q is surjective since $0, 1 \in \mathbb{Z}$ map to $[0], [1]$ |
| | (iv) $\ker(q) = \text{im}(i)$... |

If $q(a) = 0$ then $a = 2k$ for some $k \in \mathbb{Z}$, i.e. $a \in 2\mathbb{Z}$, so $\ker(q) \subseteq \text{im}(i)$.

If $a \in \text{im}(i)$ then $a = 2k$ for some $k \in \mathbb{Z}$ so $q(a) = q(2k) = 0$, hence $\text{im}(i) \subseteq \ker(q)$. □

Note $\mathbb{Z} \not\cong 2\mathbb{Z} \times \mathbb{Z}/2$ e.g. the RHS has non-id elmnt $(0, 1)$ of finite order
but we can understand elmnts of \mathbb{Z} as "2 \mathbb{Z} part" plus "1 or 0".

2. $0 \rightarrow \mathbb{Z}/2 \xrightarrow{i} \mathbb{Z}/2 \times \mathbb{Z}/2 \xrightarrow{\pi_2} \mathbb{Z}/2 \rightarrow 0$
- | |
|--------------------|
| $a \mapsto (a, 0)$ |
| $(a, b) \mapsto b$ |

- (ii) i is injective since if $(a, 0) = (a', 0)$ then $a = a'$.
(iii) π_2 is surjective since for any $b \in \mathbb{Z}/2$, $\pi_2(0, b) = b$.
(iv) $\text{im}(i) = \ker(\pi_2)$:

$$\begin{aligned}\ker(\pi_2) &= \{(a, b) \mid 0 = q(a, b) = b\} \\ &= \{(a, 0) \mid a \in \mathbb{Z}/2\}\end{aligned}$$

and $\text{im}(i) = \{(a, 0) \mid a \in \mathbb{Z}/2\}$, so they're equal.

Note $\mathbb{Z}/2 \times \mathbb{Z}/2 \cong \mathbb{Z}/2 \times \mathbb{Z}/2$ duh...

Idea: A SES $e \rightarrow N \rightarrow G \rightarrow Q \rightarrow e$ is split "if $G \cong N \times Q$ ". Otherwise, non-split.

Claim $e \rightarrow N \xrightarrow{i} N \times Q \xrightarrow{p} Q \rightarrow e$ is a SES

$n \mapsto (n, e)$
$(n, g) \mapsto g$

- (iv) $\text{im}(i) = \ker(p)$

(\Leftarrow) Suppose $(n, g) \in \text{im}(i)$. Then $g = e$, so $p(n, g) = p(n, e) = e \in Q$ so $(n, g) \in \ker(p)$.

(\Rightarrow) Suppose $(n, g) \in \ker(p)$, so $p(n, g) = e$. This means $g = e$ so $(n, g) = (n, e) = i(n) \in \text{im}(i)$.

□

Defn A SES $e \rightarrow N \xrightarrow{i} G \xrightarrow{p} Q \rightarrow e$ is split if $\exists r: G \rightarrow N$ s.t. $r \circ i = \text{id}_N$

Q. Why is this the same as " $G \cong N \times Q$ "?

↑ "retraction"

or "G is split extension"

A. Exc 2(2)

Prop - A SES splits iff $\exists d: G \xrightarrow{\cong} N \times Q$ s.t. TFDC:

$$\begin{array}{c} e \rightarrow N \xrightarrow{i} G \xrightarrow{p} Q \rightarrow e \\ \parallel \quad \parallel \quad \downarrow d \quad \parallel \quad \parallel \\ e \rightarrow N \xrightarrow{i_N} N \times Q \xrightarrow{\pi_2} Q \rightarrow e \end{array} \quad \begin{array}{l} \textcircled{1} \quad d(i(n)) = i_N(n) = (n, e) \\ \textcircled{2} \quad p(g) = \pi_2(d(g)) \Rightarrow d(g) = (? , p(g)) \end{array}$$

Proof: (\Leftarrow) Suppose $\exists d: G \xrightarrow{\cong} N \times Q$ which commutes appropriately. Define $r: G \rightarrow N$ by

$$\begin{aligned} r: G &\xrightarrow{d} N \times Q \xrightarrow{\pi_1} N \\ g &\mapsto d(g) \\ (n, q) &\mapsto n \end{aligned} \quad \text{i.e. } r = \pi_1 \circ d$$

Then r is a gp hom since both d and π_1 are, and $r \circ i(n) = \pi_1 \circ d(i(n)) = \pi_1(n, e) = n$.

(\Rightarrow) Suppose $\exists r: G \rightarrow N$ s.t. $r \circ i = \text{id}_N$. Define $d: G \rightarrow N \times Q$ by

$$d(g) = (r(g), p(g)).$$

Note that d is a gp hom since r and p are, and TFDC:

$$\begin{array}{ccc} \begin{array}{c} g \xrightarrow{p(g)} \\ \uparrow \alpha \downarrow \\ G \xrightarrow{d} N \times Q \\ \parallel \\ N \times Q \xrightarrow{\pi_1} N \\ (r(g), p(g)) \mapsto p(g) \end{array} & \begin{array}{c} n \xrightarrow{i(n)} \\ \parallel \\ N \xrightarrow{i_N} N \times Q \\ \parallel \\ N \xrightarrow{\pi_2} Q \\ (n, e) \end{array} & \star r(i(n)) = (r \circ i)(n) = n \\ & \downarrow d & \text{and } p(i(n)) = e \text{ since} \\ & (r(i(n)), p(i(n))) & \text{im}(i) = \ker(p). \end{array}$$

Remains to show d is bijective.

(injective). Suppose $d(g) = e$, so $r(g) = e = p(g)$. This means $g \in \ker(p) = \text{im}(i)$, so $g = i(n)$ for some $n \in N$. But then $e = r(g) = r(i(n)) = n$, so $n = e$. Hence $g = i(e) = e$.

(surjective). Let $(n, q) \in N \times Q$. WTS $\exists g$ s.t. $(n, q) = d(g) = (r(g), p(g))$

Know ASIDE

- Since p is surjective, $\exists \hat{g} \in G$ s.t. $p(\hat{g}) = q$, but probably not $r(\hat{g}) = n$.
- $d(i(n)) = (r(i(n)), p(i(n))) = (n, e)$

Idea: combine them

$$\begin{aligned} \bullet \quad d(i(n)\hat{g}) &= d(i(n))d(\hat{g}) = (n, e) \cdot (r(\hat{g}), q) = (nr(\hat{g}), q) \text{ almost...} \\ \bullet \quad d(i(n)\hat{g}i(r(\hat{g})^{-1})) &= d(i(n))d(\hat{g})d(i(r(\hat{g})^{-1})) \\ &= (n, e) \cdot (n\hat{g}, q) \cdot (r(\hat{g})^{-1}, e) \\ &= (n, r(\hat{g})r(\hat{g})^{-1}, e) = (n, q). \end{aligned}$$

Set $g = i(n) \cdot \hat{g} \cdot i(r(\hat{g})^{-1})$. Then $\mathcal{L}(i(n) \hat{g} i(r(\hat{g})^{-1})) = (nq)$ by above. \square

UPSHOT: Split extensions are when $G \cong N \times Q$ in a nice way. These always exist + are nice, but sometimes more interesting things happen.

Ex. two extensions of $\mathbb{Z}/2$ by $\mathbb{Z}/3$

Split: $0 \rightarrow \mathbb{Z}/3 \rightarrow \mathbb{Z}/6 \rightarrow \mathbb{Z}/2 \rightarrow 0$

$$\downarrow \text{NS} \quad \uparrow$$

$$\mathbb{Z}/3 \times \mathbb{Z}/2$$

non-split: $0 \rightarrow \mathbb{Z}/3 \xrightarrow{i} S_3 \xrightarrow{p} \mathbb{Z}/2 \rightarrow 0$
b/c $S_3 \not\cong \mathbb{Z}/6$

$$0 \mapsto e$$

$$1 \mapsto (123)$$

$$2 \mapsto (123)^2 = (132)$$

What is p ? $S_3 / i(\mathbb{Z}/3) = \{e, \langle(123)\rangle, \langle(12)\rangle \langle(123)\rangle\}$

$$\cong \downarrow$$

$$\mathbb{Z}/2$$

$$0$$

$$1$$

So $p: S_3 \rightarrow \mathbb{Z}/2$ from univ. prop.

$$\begin{array}{ccc} \downarrow & \nearrow \cong & \left. \begin{array}{c} e \\ (123) \\ (132) \\ (12) \\ (13) \\ (23) \end{array} \right\} \rightarrow 0 \\ S_3 / i(\mathbb{Z}/3) & & \left. \begin{array}{c} e \\ (123) \\ (132) \\ (12) \\ (13) \\ (23) \end{array} \right\} \rightarrow 1 \end{array}$$

Note $\mathbb{Z}/3 \cong A_3$ so similar argument shows $0 \rightarrow A_3 \hookrightarrow S_3 \rightarrow \mathbb{Z}/2 \rightarrow 0$ is non-split exact.

Special examples: Group Presentations \Leftarrow special kind of f.g. group

One of the ways we describe groups is using generators and relations

e.g. $\mathbb{Z}/n \cong \{[0], [1], \dots, [n-1]\} = \langle 1 \mid n \cdot 1 = 0 \rangle$

$D_4 = \{e, r, r^2, r^3, sr, sr^2, sr^3\}$ s.t. $r^4 = e, rs = sr^{-1}$
 $= \langle r, s \mid r^4 = e, rs = sr^{-1} \rangle$

$\mathbb{Z} = \langle 1 \rangle = \langle 1 \mid \phi \rangle$.

In general, $G = \langle \text{letters} \mid \text{relations} \rangle$ and the elmts of G are words.

e.g. $1+2+4-3$ is a word in $\mathbb{Z}/3$

$= 4-3=1 \equiv 1$ using relns \hookleftarrow

$rsr^3ss^{-1}r^2$ is a word in D_4 "reduced words"

$= rsr^3er^2 = rsr^5 = rsr$ using relns \hookleftarrow

If $\{\text{relations}\} = \emptyset$, G is called a free group

Defn For any set S , the free group on S is the group $F(S) = \langle \{seS \mid \emptyset\} \rangle$.

Intuition S is like a "basis" for $F(S)$.

This intuition is bad b/c (1) not every gp "has a basis" (i.e. is free)

(2) this "basis" is non-commutative

e.g. finite gps

Better intuition(?) S is an "alphabet" for $F(S)$.

Rmk. • $F(\emptyset) = e$

• The rank of $F(S)$ is $|S|$, and $F(S) \cong F(S') \iff |S|=|S'|$.

• UP: $\text{Hom}_{\text{Grp}}(F(S), G) \cong \text{Hom}_{\text{Set}}(S, G)$ as a set

• $F(\{s\}) \cong \mathbb{Z}$

* • $F(S)$ is non-Abelian iff $|S| \geq 2$.

• If $H \subseteq F(S)$ then H is free (i.e. $H \cong F(S')$ for some S')

e.g. $F(\{x, y\}) \not\cong F(\{x\}) \times F(\{y\}) \cong \mathbb{Z} \times \mathbb{Z}$

$$\mathbb{Z} * \mathbb{Z} \xrightarrow{\psi} x^{-3} y^2 x x y^2 y^{-1} x \quad (x^k, y^j)$$

Q. How do I get from $F(\{x, y\})$ to $F(\{x\}) \times F(\{y\})$?

A. "Abelianize": impose reln $xy = yx \iff xy(yx^{-1}) = e$

\Rightarrow quotient by "commutator subgp"

$$xyx^{-1}y^{-1} = [x, y]$$

$$F(\{x, y\}) /_{\substack{\parallel \\ [[F(\{x, y\}), F(\{x, y\})]]}} \cong F\{x\} \times F\{y\}.$$

$$\langle x, y \mid xyx^{-1}y^{-1} \rangle$$

Rmk. Can "Abelianize" any G by declaring $gh = hg$ in this way

$$G^{ab} := G / [G, G].$$

Abstract Observation if $G = \langle S \mid R \rangle$ then " $G \cong F(S)/R$ "

i.e. ESES $1 \rightarrow N(R) \rightarrow F(S) \rightarrow G \rightarrow 1$

Smallest normal subgp containing $R \subseteq F(S)$.

Ex. We saw $\langle x, y \mid xyx^{-1}y^{-1} \rangle \cong \mathbb{Z}^2$

$$\text{but } \langle x, y \mid x^4, y^2, xyx^{-1}y \rangle \cong D_4$$

Ex (Exc3) $G = \langle x, y \mid xy^2 = y^2x, x^4 = y^3 \rangle$. Show: all elmts of G can be written $y^r x^s$.

From (1), even powers of y commute w/ x + can be moved to the front. From (2), $y = y^2 x^4$
So every word in G can be expressed using only even powers of y . Hence of the form $y^r x^s$

Then: Show G Ab

$$\text{Show } G \cong \mathbb{Z}$$

Ex. $G = \langle x, y \mid xyx^{-1}y^{-2}, x^{-2}y^{-1}xy \rangle$

Note $xyx^{-1}y^{-2} = e \iff xy = (x^{-1}y^{-2})^{-1} = y^2x$

and $x^{-2}y^{-1}xy = e \iff xy = (x^2y^2)^{-1} = yx^2$

So $y^2x = yx^2$ in G . Multiplying by y^{-1} , we get $yx = x^2$ so multiplying by x^{-1} we get $y = x$. Then $xy = y^2x$ implies $x^2 = x^3$ so $x = e$ and $y = e$.

Thus $G = e$!

Hard Word Problems (early 1900s)

(1) Word Problem : can we decide when 2 words are equal?

(2) Isomorphism problem : can we decide when $\langle S | R \rangle \cong \langle S' | R' \rangle$?

↳ e.g. $\langle S | R \rangle = e$?

A. It's undecidable.

~ break ~

HW time : Exercise 4 only!

Exercise 4. We show in this exercise that every group is determined by its finitely generated subgroups. Conceptually, this is not surprising. The multiplication gh in G for $g, h \in G$ is entirely determined in the finitely generated subgroup $\langle g, h \rangle$.

- (1) A *partially ordered set* $I = (I, \leq)$ is a non-empty set I together with a relation \leq which is reflexive, antisymmetric (i.e. if $a \leq b$ and $b \leq a$, then $a = b$), and transitive. A *filtered set* $I = (I, \leq)$ is a partially ordered set together with upper bounds: for all $a, b \in I$, there exists $c \in I$ such that $a \leq c$ and $b \leq c$. Show that (\mathbb{N}, \leq) and (\mathbb{R}, \leq) are filtered sets.
- (2) A *filtered system of groups* $I = (I, \leq)$ is a filtered set in groups: each element in I is a group and we fix injections $\iota_{HK} : H \rightarrow K$ for some pairs of groups $H, K \in I$. The relation \leq is defined as $H \leq G$ if and only if we have chosen an injective homomorphism $\iota_{H,G} : H \rightarrow G$, for H, G in I . Fix now a group G and let I_G be the set of all finitely generated subgroups of G . Show that I_G is a filtered system of groups.
- (3) Let I be a filtered set of groups. We define the *filtered colimit* $(G, \{f_H\}_{H \in I})$ of I as follows. It is a group G together with homomorphisms $f_H : H \rightarrow G$ for each $H \in I$ such that for all injective homomorphisms $\iota_{HK} : H \hookrightarrow K$ in the filtered system I , we have $f_K \circ \iota_{HK} = f_H$ for all $H, K \in I$. It respects a universal property that reads: for any other group G' with homomorphisms $\{f'_H\}_{H \in I}$ such that $f'_K \circ \iota_{HK} = f'_H$ for all $H, K \in I$, then there exists a unique homomorphism $F : G \rightarrow G'$ such that $F \circ f_H = f'_H$ for all $H \in I$. Show that given a filtered set of groups I , the filtered colimit of I is unique up to isomorphism if it exists.
- (4) Let G be a group and let I_G be the filtered set as in (2). Show that G is the filtered colimit of I_G .

minute sheet

- which isomorphism is the best iso form?
- what was most helpful today?