

**Correction to the proof of Lemma 9.4.2 in  
*Repeated Games and Reputations*  
 by George J. Mailath and Larry Samuelson (OUP 2006)  
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Claim 9.4.3 (on page 308) is false. We thank Songzi Du for bringing the error in the proof to our attention and providing a counterexample to the claim. There does not appear to be an easy way to patch the proof presented in the text of Lemma 9.4.2 (based on Fudenberg, Levine, and Maskin (1994)). The following is an alternative proof, also due to Fudenberg, Levine, and Maskin.

**Lemma 9.4.2** *If an action profile with pure long-lived players actions is enforceable and pairwise identifiable for long-lived players  $i$  and  $j$ , then it is orthogonally enforceable in all  $ij$ -pairwise directions,  $\lambda^{ij}$ .*

**Proof.** Suppose  $\alpha^* = (a_1^*, \dots, a_n^*, \alpha_{SL}^*)$  is an action profile with pure long-lived players actions enforced by some normalized continuations  $\hat{x}$ . Let  $v_i = u_i(\alpha^*) + E[\hat{x}_i(y) \mid \alpha^*]$ , and define  $g_i(a_i) \equiv u_i(a_i, \alpha^*) - v_i$  for all  $a_i \in A_i$ . We first prove a preliminary claim.

**Claim 1** *For all  $i$ , and all  $\mu_i \in \mathbb{R}^{A_i}$  satisfying*

$$\sum_{a_i \in A_i} \mu_i(a_i) \rho(y \mid (a_i, \alpha_{-i}^*)) = 0, \quad \forall y \quad (1)$$

$$\text{and} \quad \mu_i(a_i) \geq 0, \quad \forall a_i \neq a_i^*, \quad (2)$$

*we have*

$$\sum_{a_i \in A_i} \mu_i(a_i) g_i(a_i) \leq 0. \quad (3)$$

**Proof.** By hypothesis, the normalized continuations  $\hat{x}_i : Y \rightarrow \mathbb{R}$  satisfy (8.1.2) and (8.1.3) at  $\alpha^*$ . We rewrite (8.1.2) and (8.1.3) as

$$\sum_{y \in Y} \rho(y \mid (a_i^*, \alpha_{-i}^*)) x_i(y) = -g_i(a_i^*), \quad (4)$$

$$\text{and} \quad \sum_{y \in Y} \rho(y \mid (a_i, \alpha_{-i}^*)) x_i(y) \leq -g_i(a_i), \quad \forall a_i \in A_i \setminus \{a_i^*\}. \quad (5)$$

These linear constraints can be viewed as the constraints of a linear program, with an objective function identically equal to 0. That is,

since  $\alpha^*$  is enforceable, the following linear program has a solution (for all  $i$ ):

$$\max_{x_i \in \mathbb{R}^Y} \sum_{y \in Y} 0 \cdot x_i(y) \quad \text{subject to (4) and (5)}. \quad (\text{LP}_i)$$

Since the continuations are not constrained to be nonnegative, the dual of this linear program is

$$\max_{\mu_i \in \mathbb{R}^{A_i}} \sum_{a_i \in A_i} \mu_i(a_i) g_i(a_i) \quad \text{subject to}$$

$$\sum_{a_i \in A_i} \mu_i(a_i) \rho(y | (a_i, \alpha_{-i}^*)) = 0, \quad \forall y \quad (6)$$

$$\text{and } \mu_i(a_i) \geq 0, \quad \forall a_i \neq a_i^*. \quad (7)$$

Note that (6) coincides with (1), and (7) coincides with (2). The dual has at least one feasible vector of dual variables (given by  $\mu_i(a_i) = 0$  for all  $a_i$ ), and so the primal program  $\text{LP}_i$  has a solution if, and only if, the dual is bounded.<sup>1</sup> Since the primal program has a solution, the dual is bounded.

Suppose  $\mu_i \neq 0$  is a feasible vector of dual variables. Then,  $\mu_i(a_i^*) < 0$  (if not, the dual constraints imply  $\rho(y | (a_i, \alpha_{-i}^*)) = 0$  for all  $y$  and any  $a_i$  satisfying  $\mu_i(a_i) > 0$ , which is impossible). Rewrite (6) as (where  $B_i \equiv \{a_i \in A_i : \mu_i(a_i) \geq 0\} = A_i \setminus \{a_i^*\}$ ),

$$\sum_{a_i \in B_i} \mu_i(a_i) \rho(y | (a_i, \alpha_{-i}^*)) = -\mu_i(a_i^*) \rho(y | \alpha^*). \quad (8)$$

Since  $\mu_i \neq 0$ , we have  $\sum_{a_i \in B_i} \mu_i(a_i) > 0$ . Defining  $\tilde{\mu}_i(a_i) \equiv \mu_i(a_i) / \sum_{a_i \in B_i} \mu_i(a_i)$ , we have a well-defined probability distribution over actions in  $B_i$ . Moreover,

$$\begin{aligned} 1 &= \sum_y \sum_{a_i \in B_i} \tilde{\mu}_i(a_i) \rho(y | (a_i, \alpha^*)) \\ &= - \sum_y \tilde{\mu}_i(a_i^*) \rho(y | \alpha^*) = -\tilde{\mu}_i(a_i^*). \end{aligned}$$

Equation (8) thus implies that the mixed action induced by  $\tilde{\mu}_i$  on  $B_i$  implies the *same* distribution on the public signal as  $a_i^*$ .

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<sup>1</sup>For a compact introduction to linear programming, see Vohra (2005).

In other words, for any feasible nonzero vector  $\mu_i$ , the objective function can be written as

$$\sum_{a_i \in B_i} \mu_i(a_i) \left\{ \sum_{a_i \in B_i} \tilde{\mu}_i(a_i) g_i(a_i) - g_i(a_i^*) \right\}.$$

Note that if  $\mu_i$  is a feasible vector, then so is  $k\mu_i$  for any positive  $k$ . The requirement that the dual be bounded over the set of feasible dual variables is hence equivalent to the requirement that for all dual variables satisfying (6) and (7), we have (3). ■

The action profile  $\alpha^*$  is orthogonally enforceable in the pairwise direction  $\lambda^{ij}$  if there exist normalized continuations  $x: Y \rightarrow \mathbb{R}^n$  satisfying (4) and (5) for all long-lived players, and  $\lambda_i^{ij} x_i(y) + \lambda_j^{ij} x_j(y) = 0$  for all  $y \in Y$ . As usual, since the pairwise direction  $\lambda^{ij}$  imposes no further restrictions on the normalized continuations for the long-lived players other than  $i$  and  $j$ , we may ignore them.

We now argue that the normalized continuations  $\hat{x}_i$  and  $\hat{x}_j$  that enforce  $\alpha^*$  (and yield payoffs  $v_i$  and  $v_j$ ) can be adjusted to obtain orthogonal enforceability without affecting the values  $v_i$  and  $v_j$ . As in the proof of the claim, the existences of the desired normalized continuations for  $i$  and  $j$  is equivalent to the following linear program having a solution (where we have eliminated  $x_j$  using the orthogonality condition):

$$\begin{aligned} \max_{x_i \in \mathbb{R}^Y} \quad & \sum_{y \in Y} 0 \cdot x_i(y) \quad \text{subject to} \\ & \sum_{y \in Y} \rho(y | (a_i^*, \alpha_{-i}^*)) x_i(y) = -g_i(a_i^*), \\ & \sum_{y \in Y} \rho(y | (a_i, \alpha_{-i}^*)) x_i(y) \leq -g_i(a_i), \quad \forall a_i \in A_i \setminus \{a_i^*\}, \\ & \sum_{y \in Y} \rho(y | (a_j^*, \alpha_{-j}^*)) (-\lambda_i^{ij} / \lambda_j^{ij}) x_i(y) = -g_j(a_j^*), \\ \text{and} \quad & \sum_{y \in Y} \rho(y | (a_j, \alpha_{-j}^*)) (-\lambda_i^{ij} / \lambda_j^{ij}) x_i(y) \leq -g_j(a_j), \quad \forall a_j \in A_j \setminus \{a_j^*\}. \end{aligned}$$

The dual of this linear program is

$$\max_{\mu_i \in \mathbb{R}^{A_i}, \mu_j \in \mathbb{R}^{A_j}} \sum_{a_i \in A_i} \mu_i(a_i) g_i(a_i) + \sum_{a_j \in A_j} \mu_j(a_j) g_j(a_j)$$

subject to

$$\sum_{a_i \in A_i} \mu_i(a_i) \rho(y | (a_i, \alpha_{-i}^*)) + \mu_j(a_j) \rho(y | (a_j, \alpha_{-j}^*)) (-\lambda_i^{ij} / \lambda_j^{ij}) = 0, \quad \forall y, \quad (9)$$

$$\mu_i(a_i) \geq 0, \quad \forall a_i \neq a_i^*, \quad (10)$$

$$\text{and } \mu_j(a_j) \geq 0, \quad \forall a_j \neq a_j^*. \quad (11)$$

As in the proof of the claim, there exist orthogonally enforcing continuations if the dual problem is bounded for feasible dual variables. We now argue that this pairwise problem can be effectively separated into distinct individual problems.

The system (9) can be written as

$$(\mu_i, \hat{\mu}_j) \cdot R_{ij}(\alpha^*) = 0, \quad (12)$$

where  $\hat{\mu}_j = -\lambda_i^{ij} \mu_j / \lambda_j^{ij}$ . Note that  $\hat{\mu}_j$  satisfies (1) if, and only if,  $\mu_j$  does.

For an arbitrary matrix  $R$ , denote its “left” null space by  $N(R)$ , that is,  $N(R) = \{\xi : \xi \cdot R = 0\}$ . Thus, (1) is the requirement that  $\mu_i \in N(R_i(\alpha_{-i}^*))$ , while (12) is the requirement that  $(\mu_i, \hat{\mu}_j) \in N(R_{ij}(\alpha^*))$ . Let  $\dim(N(R))$  denote the dimension of the left null space of  $R$ . Pairwise identifiability implies

$$\dim(N(R_{ij}(\alpha^*))) = \dim(N(R_i(\alpha_{-i}^*))) + \dim(N(R_j(\alpha_{-j}^*))) + 1.$$

Since  $(\alpha_i^*, -\alpha_j^*) \in N(R_{ij}(\alpha^*))$  (recall the comment just after Definition 9.2.2), every vector in  $N(R_{ij}(\alpha^*))$  is a linear combination of  $(\alpha_i^*, -\alpha_j^*)$  and some  $(\mu_i, \mu_j)$ , where  $\mu_i \in N(R_i(\alpha_{-i}^*))$  and  $\mu_j \in N(R_j(\alpha_{-j}^*))$ . Note that the value of the objective function at  $(\alpha_i^*, -\alpha_j^*)$  is zero (since  $g_i(a_i^*) = g_j(a_j^*) = 0$ ), and so the weight placed on that vector can be ignored.

Since (10) and (11) are each (2) for players  $i$  and  $j$ , the claim implies that for any  $(\mu_i, \mu_j)$ , where  $\mu_i \in N(R_i(\alpha_{-i}^*))$  and  $\mu_j \in N(R_j(\alpha_{-j}^*))$ , (3) holds for  $i$  and  $j$ , i.e., the dual for the pairwise problem is bounded above by 0. Hence, the primal for the pairwise problem has a solution and so  $\alpha^*$  is orthogonally enforceable in all  $ij$ -pairwise directions. ■

## References

- FUDENBERG, D., D. K. LEVINE, AND E. MASKIN (1994): “The Folk Theorem with Imperfect Public Information,” *Econometrica*, 62(5), 997–1039.
- VOHRA, R. V. (2005): *Advanced Mathematical Economics*. Routledge, London and New York.